

MATHEMATICAL REASONING

EXERCISE - 01

CHECK YOUR GRASP

5. $\therefore (p \rightarrow \sim p) \wedge (\sim p \rightarrow p)$
 $\equiv (\sim p \vee \sim p) \wedge (p \vee p) \quad (\because p \rightarrow q \equiv \sim p \vee q)$
 $\equiv \sim p \wedge p \quad (\because p \vee p \equiv p)$
 $\equiv c$ c is a contradiction
8. $\therefore \sim q \rightarrow \sim p \equiv \sim(\sim q) \vee \sim p \quad (\because p \rightarrow q \equiv \sim p \vee q)$
 $\equiv q \vee \sim p$
 $\equiv \sim p \vee q \quad (\text{by commutative law})$
 $\equiv p \rightarrow q \quad (\because p \rightarrow q \equiv \sim p \vee q)$
Hence $p \rightarrow q \equiv \sim q \rightarrow \sim p$
12. $\therefore (p \wedge q) \rightarrow p$ is false
 $\Rightarrow (p \wedge q)$ is true and p is false
which is not possible
so $(p \wedge q) \rightarrow p$ is always true i.e. it is a tautology.
14. $\therefore p \rightarrow (q \vee r)$ is false
 $\Rightarrow p$ is true and $(q \vee r)$ is false
 $\Rightarrow p$ is true, q and r both are false
i.e. $p \rightarrow (q \vee r)$ is false when truth values of p, q, r are T, F, F resp. otherwise it is true.
15. $\therefore (p \wedge q) \wedge \sim(p \vee q)$ is true
 $\Rightarrow (p \wedge q)$ and $\sim(p \vee q)$ both are true
which is not possible
So $(p \wedge q) \wedge \sim(p \vee q)$ is always false i.e. it is a contradiction.
16. Let p, q, r three statement defined as
 p : a number N is divisible by 15
 q : number N is divisible by 5
 r : number N is divisible by 3
Here given statement is $p \rightarrow (q \vee r)$
Here negative of above statement is
 $\sim(p \rightarrow (q \vee r)) \equiv p \wedge (\sim(q \vee r))$
 $\equiv p \wedge (\sim q \wedge \sim r)$
i.e. A number is divisible by 15 and it is not divisible by 5 and 3.
19. $\therefore (p \wedge q) \vee (q \wedge r)$ is false
 $\Rightarrow (p \wedge q)$ and $(q \wedge r)$ both are false
 $\Rightarrow p$ and r both are false or q is false.
otherwise $(p \wedge q) \vee (q \wedge r)$ is true
20. $\therefore (\sim p \vee \sim q) \vee (p \vee \sim q)$
 $\equiv (\sim p \vee \sim q) \vee (\sim q \vee p) \quad (\text{by commutative law})$
 $\equiv \sim p \vee [\sim q \vee (\sim q \vee p)] \quad (\text{by Associative law})$
 $\equiv \sim p \vee [(\sim q \vee \sim q) \vee p] \quad (\text{by Associative law})$
 $\equiv \sim p \vee (\sim q \vee p) \quad (\because p \vee p \equiv p)$
 $\equiv \sim p \vee (p \vee \sim q) \quad (\text{by commutative law})$
 $\equiv (\sim p \vee p) \vee \sim q \quad (\text{by Associative law})$
 $\equiv t \vee \sim q \equiv t$ t is a tautology
Hence $(\sim p \vee \sim q) \vee (p \vee \sim q)$ is a tautology.
24. When p and q both are true then
 $\sim(p \rightarrow q)$ and $(\sim p \vee \sim q)$ both are false
i.e. $\sim(p \rightarrow q) \leftrightarrow (\sim p \vee \sim q)$ is true
when p and q both are false then
 $\sim(p \rightarrow q)$ is false and $(\sim p \vee \sim q)$ is true
i.e. $\sim(p \rightarrow q) \leftrightarrow (\sim p \vee \sim q)$ is false
Hence $\sim(p \rightarrow q) \leftrightarrow (\sim p \vee \sim q)$ is neither tautology nor contradiction.
28. Let $S(p, q) \equiv (p \vee \sim q) \wedge \sim p$
 $\Rightarrow S(\sim p, \sim q) \equiv (\sim p \vee q) \wedge p$
Now $S^*(\sim p, \sim q) \equiv (\sim p \wedge q) \vee p$
and $\sim S(p, q) \equiv \sim[(p \vee \sim q) \wedge \sim p] \equiv \sim(p \vee \sim q) \vee p$
 $\equiv (\sim p \wedge q) \vee p$
Hence $S^*(\sim p, \sim q) \equiv \sim S(p, q)$
32. $\therefore [(p \wedge p) \rightarrow q] \rightarrow p \equiv (p \rightarrow q) \rightarrow p \quad (\because p \wedge p \equiv p)$
when p is false and q is true (or false) then
 $(p \rightarrow q)$ is true i.e. $(p \rightarrow q) \rightarrow p$ is false
Hence $[(p \wedge p) \rightarrow q] \rightarrow p$ is not a tautology.
34. $\therefore (p \vee \sim r) \rightarrow (q \wedge r)$ is false
 $\Rightarrow (p \vee \sim r)$ is true, and $(q \wedge r)$ is false
Here $(q \wedge r)$ is false and q is true (given)
 $\Rightarrow r$ is false
again r is false and $(p \vee \sim r)$ is true
 $\Rightarrow p$ may be true or false.

EXERCISE - 02

1. $\therefore p \rightarrow (q \rightarrow p)$ is false
 $\Rightarrow p$ is true and $(q \rightarrow p)$ is false. which is not possible.
 So $p \rightarrow (q \rightarrow p)$ is always true i.e. it is a tautology.
 Again $p \rightarrow (p \vee q)$ is false
 p is true and $(p \vee q)$ is false. Which is not possible.
 So $p \rightarrow (p \vee q)$ is always true i.e. it is a tautology.
 Hence $p \rightarrow (q \rightarrow p) \equiv p \rightarrow (p \vee q)$
2. Given $r : \sim p \leftrightarrow q$
 statement-1 $r \equiv q \vee p$
 statement-2 $r \equiv (p \leftrightarrow \sim q)$

p	q	$\sim p$	$\sim q$	$(\sim p \leftrightarrow q)$	$q \vee p$	$(p \leftrightarrow \sim q)$
T	T	F	F	F	T	F
T	F	F	T	T	T	T
F	T	T	F	T	T	T
F	F	T	T	F	F	F

Hence Statement-1 is false and Statement-2 is true.

3. statement-1 : $\sim(p \leftrightarrow \sim q)$ is equivalent to $p \leftrightarrow q$
 statement-2 : $\sim(p \leftrightarrow \sim q)$ is a tautology.

p	q	$\sim q$	$(p \leftrightarrow q)$	$(p \leftrightarrow \sim q)$	$\sim(p \leftrightarrow \sim q)$
T	T	F	T	F	T
T	F	T	F	T	F
F	T	F	F	T	F
F	F	T	T	F	T

Hence statement-1 is true, statement-2 is false.

PREVIOUS YEAR QUESTION

4. Given $S \subseteq R$ and
 $p = \text{There is a rational number } x \in S \text{ such that } x > 0$
 then $\sim p : \text{Any rational number } x \in S \text{ such that } x \not> 0$
 i.e. $\sim p : \text{Every rational number } x \in S \text{ satisfy } x \leq 0$
5. Given Statement :
 $(p \wedge \sim r) \leftrightarrow q$
 Negations of $p \leftrightarrow q$ are
 $\sim(p \leftrightarrow q), \sim(q \leftrightarrow p),$
 $\sim p \leftrightarrow q$ and $\sim q \leftrightarrow p$
 Hence negations of given statement
 are $\sim(q \leftrightarrow (p \wedge \sim r))$
 and $\sim(p \wedge \sim r) \leftrightarrow q$
6. $[p \wedge (p \rightarrow q)] \rightarrow q$
 $[p \wedge (\sim p \vee q)] \rightarrow q$
 $[(p \wedge \sim p) \vee (p \wedge q)] \rightarrow q$
 $[c \vee (p \wedge q)] \rightarrow q \quad \begin{cases} p \wedge \sim p \equiv c \equiv \text{contradiction} \\ \therefore c \vee p \equiv p \end{cases}$
 $\Rightarrow (p \wedge q) \rightarrow q$
 $\Rightarrow \sim(p \wedge q) \vee q$
 $\Rightarrow (\sim p \vee \sim q) \vee q$
 $\Rightarrow \sim p \vee (q \vee \sim q)$
 $\Rightarrow \sim p \vee (t) \equiv \text{tautology}$